A Direct Non-symmetry and Anti-packing Model for Color Images

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Abstract

The representation methods of the hierarchical data structures have been widely applied in many important fields, such as computer visualization, robotics, computer graphics, image processing, and pattern recognition. Since these methods put too much emphasis upon the symmetry of segmentation, they are not the optimal representation methods. In this paper, we present a Direct Non-symmetry and Anti-packing Model (DNAM) for color images. Also, we propose a novel algorithm of the DNAM for color images and analyze the data amount of the proposed algorithm. By taking a rectangle subpattern for example, we implement the proposed algorithm and make a comparison with the algorithm of the popular linear quadtree. The theoretical and experimental results presented in this paper show that our proposed algorithm can reduce the data storage much more effectively than that of the linear quadtree and it is a better method to represent color images.

1. Introduction

The hierarchical quadtree-based representation methods have been widely applied in computer visualization, robotics, computer graphics, image processing, and pattern recognition. Klinger [1] proposed a quadtree representation method for a binary image. However, the earlier quadtree representation is based on the pointers. Later, in order to reduce the storage room further, Gargantini [2] removed the pointers and put forward a linear quadtree representation method. The establishment of the linear quadtree representation is an important milestone for the history of the quadtree representation. Until recently, the linear quadtree method is still the most popular representation and is widely studied and applied in many fields [3-8]. However, although the representation methods of the hierarchical data structures stated above have many merits and applications, they put too much emphasis upon the symmetry of segmentation. Therefore, they are not the optimal representation methods. Inspired by the concept

of the rectangle packing problem, Zheng and Chen presented a color image representation method based on Non-symmetry and Anti-packing Model (NAM) by using the method of Binary-bit Plane Decomposition (BPD) in [9] where the predefined subpatterns are rectangles, since the traditional representation methods of the hierarchical data structures have their limitations in the representation efficiency. Also, a novel algorithm was put forward for the triangle non-symmetry and anti-packing pattern representation model (TNAM) of gray images in [10] where the predefined subpatterns are triangles.

In this paper, without using the method of BPD [9] and the triangle subpatterns [10], we present a Direct Nonsymmetry and Anti-packing pattern representation Model (DNAM) [11] for color images. Also, we propose a novel algorithm of the DNAM for color images and analyze the data amount of this algorithm. The theoretical and experimental results presented in this paper show that our proposed algorithm of the DNAM can reduce the data storage much more effectively than that of the popular linear quadtree and it is a better method to represent color images.

2. The idea of Direct Non-symmetry and Anti-packing Model (DNAM)

A detail description about the DNAM can refer to [11]. In this paper, the idea of the DNAM for color images can be described as following: Giving a packed color image pattern and n predefined subpatterns with different shapes, pick up these subpatterns from the packed pattern and then represent the pattern with the combination of these subpatterns.

We know that a color image can be represented by using three gray images. By using the method of BPD [9], a gray image with a gray level m can be decomposed to m binary images. The method of BPD can effectively reduce the complexity of a gray image, and can make the total subpatterns number of all of the decomposed binary images much less than that of the gray image.

However, in this paper, as far as the DNAM for color images with a gray level m is concerned, we do not treat

each gray image as m binary images. On the contrary, we treat each gray image just as an indecomposable gray image.

3. A novel algorithm of the DNAM for color images

3.1. Rules of K-Code transform

Before describing the encoding and decoding algorithm of the DNAM for color images, we first reintroduce the rules of K-Code transform [9,10] since our algorithm strongly depends on them.

Suppose a binary image of size $2^n \times 2^n$ is denoted by $F = \{f(x,y)\}$, i.e., $f(x,y) \{0,1\}$. Without loss of generality, we assume that 0 and 1 stand for the black and white pixels, respectively, and that the white pixels denote the backgrounds, whereas the black pixels denote the objects. Let (x,y) be a two-dimensional coordinate, where $x=(x_n, 1x_{n-2}...x_1x_0)_2$ and $y=(y_{n-1}y_{n-2}...y_1y_0)_2$. By rearranging the binary bits of x and y, we can obtain a new coordinate variable k, i.e., $k=(y_{n-1}x_{n-2}...y_1x_0x_0)_2$. Herein, the two-dimensional image is converted into a one-dimensional sequence, i.e., $F=\{f(x,y)\}=\{g(k)\}$. This one-dimensional representation for the image is called as K-Code.

The transform from a two-dimensional coordinate (x,y) to a one-dimensional coordinate k is called as K-Code transform, i.e., k=K(x,y). On the contrary, the transform from k to (x,y) is called as Counter-K-Code transform, i.e., $(x,y)=K^{-1}(k)$.

3.2. Description of our proposed algorithm of the DNAM for color images

In this section, by taking a single rectangle subpattern $p=\{\text{rectangle}|\text{rectangle}=L\times W\}$ for example, we propose a novel algorithm of the DNAM for color images. The following procedures Rect_encoder () and Rect_decoder () describe the encoding and decoding parts of the proposed algorithm of the DNAM for color images, respectively. The K-Code transform rules are adopted to reduce the total subpatterns number.

procedure Rect_encoder (n, m, {CI(x,y)}, Q) /* Given a color image {CI(x,y)} of size $2^n \times 2^n \times 3$ with a gray level m and a resolution n, an encoding result is stored into a queue Q. */ **begin**

color image {CI(x,y)}; gray image {MP[1](x,y)}, {MP[2](x,y)}; gray image {MP[3](x,y)}; K-code $k_0=0, \{g(k)\}; /*$ K-Code array of size $4^n */$ DNAM table $Q=\{Q[1], Q[2], Q[3]\};$ integer n, m, gray, MAXarea, area, j, k, x, y, t; integer i, j, length, width, length₀, width₀;

for t \leftarrow 1 to 3 begin $MP[t] \leftarrow rgb2gray(CI, t);$ /* Decompose a color image CI to its three gray images by three-color components r, g, and b, i.e. MP[1], MP[2], and MP[3], respectively. */ i=0; while $(i < 4^n)$ begin $(\mathbf{x},\mathbf{y}) \leftarrow \mathbf{K}^{-1}(\mathbf{i});$ if MP[t](x,y) \neq -1 then begin gray \leftarrow MP[t](x,y); length $\leftarrow 1$; width $\leftarrow 1$; MAXarea $\leftarrow 1$; for $i \leftarrow y$ to $2^n - 1$ do begin for k \leftarrow x to 2ⁿ-1 do begin if MP[t](k,j) \neq gray then break; $length_0 \leftarrow length_0 + 1;$ end area \leftarrow length₀ × width₀; if MAXarea<area then begin MAXarea←area; length \leftarrow length₀; width \leftarrow width₀; end if MP[t](x,j) \neq gray then break; width₀←width₀+1 end $Q[t] \leftarrow \{gray, length, width\};$ marking the rectangle as -1; end $i \leftarrow i+1$; end end output(Q); end procedure Rect decoder $(n, m, Q, \{CI(x,y)\})$ /* Given a queue Q, a resolution n, and a gray level m of a color image $\{CI(x,y)\}$, a decoding result is stored into {CI(x,y)}. */ begin color image $\{CI(x,y)\};$ gray image {MP[1](x,y)}= $\{-1\}$, {MP[2](x,y)}= $\{-1\}$; gray image $\{MP[3](x,y)\} = \{-1\};$ K-Code k₀=0; DNAM table $Q = \{Q[1], Q[2], Q[3]\};$ integer max, ptr, n, i, gray, length, width, x, y, j; for $i \leftarrow 1$ to 3 begin ptr=0; max←size(Q[i]); while (ptr<max) begin

```
{gray, length, width}\leftarrow Q[i](ptr);
(x,y) \leftarrow top_left_coordinate(gray);
for j\leftarrow0 to width-1 do
begin
for k\leftarrow0 to length-1 do
begin
MP[i](x+k,y+j)=gray;
end
end
ptr\leftarrowptr+1;
end
cI\leftarrowgray2rgb(MP[1], MP[2], MP[3]);
end
```

3.3. Storage structure of the DNAM for color images

In this section, a storage structure of the DNAM for a color image of size $2^n \times 2^n \times 3$ is analyzed with respect to a typical rectangle subpattern.

The output of the procedure Rect_ encoder () is a queue Q. As far as a color image CI with a gray level m is concerned, m bits are needed to represent the gray value of CI. In addition, a rectangle subpattern has two parameters, i.e., the length L and the width W. According to the definition of K-Code, the maximal lengths of both L and W are n/2 bits. Thus, the storage structure of the DNAM for color images can be denoted as shown in Fig.1. Therefore, storing a rectangle subpattern takes up m+n bits.



Figure 1. A storage structure of the DNAM

3.4. Data amount analysis of the proposed algorithm of the DNAM for color images

In this section, the representation efficiency of the DNAM for color images is fully embodied. Before analyzing the data amount of the proposed algorithm of the DNAM for a color image *CI* of size $2^n \times 2^n \times 3$ with a gray level *m*, we first introduce the following definition [10] which describes the complexity of a color image in order that the DNAM can be conveniently compared with the linear quadtree.

Definition 1. The complexity of a color image is defined as follows:

$$Cc = N_{LQT} / N_f$$
,

where N_{LQT} is the total nodes number of the color image which is represented by the method of the linear quadtree and N_f is the size of the color image.

Suppose the total pixels number of the color image CI is N_f , i.e., $N_f=3\times 4^n$. After the image is anti-packed with a

typical rectangle subpattern, suppose the total subpatterns number and the total data amount are N_r and H_r , respectively.

As far as a typical rectangle subpattern is concerned, storing a record of the DNAM needs m+n bits. Therefore, we can write the total data amount H_r as follows:

$$H_r = (m+n)N_r = (m+n)N_f Cc \frac{N_r}{N_{LQT}}.$$

$$= 3(m+n)4^n Cc \frac{N_r}{N_{LQT}}.$$
(1)

As for the method of the linear quadtree, storing a record needs 3(n-1)+2 bits [2], and takes up 3(n-1)+2+m bits for a color image since the gray level *m* needs to be stored. Suppose that the total data amount of the color image *CI* is H_{LQT} when *CI* is represented by the method of the linear quadtree. We can write H_{LOT} as follows:

$$H_{LQT} = (3n-1+m)N_{LQT} = (3n-1+m)N_fCc .$$
(2)
= 3(3n-1+m)4ⁿCc

Then, let φ denotes the ratio of H_{LQT} to H_r . We can easily deduce the following expression (4).

$$\varphi = \frac{3(3n-1+m)4^n Cc}{3(m+n)4^n Cc \frac{N_r}{N_{LQT}}} = \frac{(3n-1+m)N_{LQT}}{(m+n)N_r} \quad (3)$$

The total data amounts of the DNAM and the linear quadtree can be achieved by (1) and (2), respectively. From these two expressions we can know that the total data amounts of the two methods are both relative to the image complexity. Since the image complexities are different for different images, the representation efficiencies are also various.

The last expression (3) shows the ratio of the total data amount of the linear quadtree to that of the DNAM. From this expression, we can judge which method is better to represent the color image pattern. Since the quadtree segmentation is symmetrical, the segmentation methods suffer from a great confine. However, since the DNAM segmentation is asymmetrical, the segmentation methods are unrestricted. The purpose of the DNAM segmentation is to form the rectangle subpatterns as big as possible and yield the least subpatterns number for a packed pattern. Therefore, generally speaking, N_r is much less than N_{LQT} , i.e., $N_r < N_{LQT}$. Thus, we can easily deduce the following expression:

$$\varphi = \frac{(3n-1+m)N_{LQT}}{(m+n)N_r} > \frac{3n-1+m}{m+n} > 1$$

For example, when m=8 and n=9, φ is greater than 2. Thus, $\varphi > 2>1$.

Therefore, the theoretical analysis in this section proves that our representation method of the DNAM for color images can reduce the data storage more effectively than that of the popular linear quadtree and it is a better method to represent color images.

4. Experimental results

To verify the obtained theoretical results, some representative color images of size $2^9 \times 2^9 \times 3$ (see Fig.2), i.e., n=9, are analyzed. The gray level of these images is 8, i.e., m=8. By taking a single rectangle subpattern for example, we have implemented the proposed algorithm of the DNAM for color images with the Matlab 6.5 software and made a comparison with the algorithm of the popular linear quadtree. The following Table 1 shows the comparison of Cc, N, and φ between the linear quadtree and the DNAM. These representative images have different image complexities. The image complexity is defined according to the nodes number of the linear quadtree. Therefore, according to the value of Cc, we can know how complex an image is. In addition, the bigger the value Cc is, the more complex the image is. From the values of Cc, we can know that the image complexities of Fig.2(a) and Fig.2(b) are relatively higher, whereas the image complexities of Fig.2(e) and Fig.2(f) are relatively lower. The nodes number N of the Quadtree in Table 1 also explains this fact. In addition, the subpatterns number of the DNAM is much less than the nodes number of the Quadtree.

Also, from the values of φ , we know that as far as the six images are concerned, the total data amount of the Quadtree is 2.3145 to 6.9067 times of that of the DNAM, which verifies the theoretical obtained result, i.e., $\varphi > 2 > 1$.

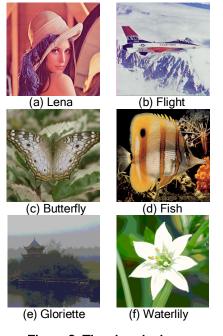


Figure 2. The six color images

Table 1. Comparison of the performance between Quadtree and DNAM

Image	Сс	N		
		Quadtree	DNAM	$- \varphi$
Lena	0.9976	677779	784545	2.3145
Flight	0.9920	630850	780159	2.4730
Butterfly	0.8895	470878	699519	2.9713
Fish	0.7847	399232	617094	3.0910
Gloriette	0.2216	62622	174309	5.5669
Waterlily	0.1953	44486	153621	6.9067

NOTE: *Cc*: image complexity; *N*: number of subpatterns or nodes; Quadtree: linear quadtree; φ : ratio of the total data amount of the Quadtree to that of the DNAM.

Therefore, the experimental results in this section show that our proposed algorithm of the DNAM for color images is much more effective than that of the popular linear quadtree with respect to the data storage and it is a better method to represent color images.

5. Conclusions and extensions

In this paper, we presented a novel Non-symmetry and Anti-packing Model (DNAM) for color images. Also, we proposed a novel algorithm of the DNAM for color images, tested the algorithm of the DNAM for color images with the Matlab 6.5 software, and compared our proposed algorithm with the popular linear quadtree's with respect to the representation efficiency. The theoretical and experimental results shown that our proposed algorithm of the DNAM for color images can reduce the data storage much more effectively than that of the popular linear quadtree.

However, although the implementation of the proposed algorithm of the DNAM for color images is based on a single rectangle subpattern, and is not sophisticated, yet it yields excellent encoding and decoding results. It seems extremely likely that well-devised improvements of this algorithm could still increase encoding performance dramatically. For example, as far as the multi-subpattern or arbitrary subpattern is concerned, we can expect a much better result than the single rectangle subpattern's.

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7. References

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