# Modified constructions of binary sequences using multiplicative inverse 

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#### Abstract

Two new families of finite binary sequences are constructed using multiplicative inverse. The sequences are shown to have strong pseudorandom properties by using some estimates of certain exponential sums over finite fields. The constructions can be implemented fast since multiplicative inverse over finite fields can be computed in polynomial time.


## §1 Introduction

In the last decade, a new constructive approach has been developed to study pseudorandomness of finite binary sequences. To the best of our knowledge, the starting work is [1]. The work was motivated by the facts that pseudorandom binary sequences have many applications such as in stream ciphers and the theory of pseudorandomness can be used to analyze certain sequences. Mauduit and Sárközy ${ }^{[1]}$ first introduced the following measures of pseudorandomness for a finite binary sequence with length $N$ :

$$
S_{N}=\left\{s_{1}, s_{2}, \cdots, s_{N}\right\} \in\{+1,-1\}^{N}
$$

The well-distribution measure of $S_{N}$ is defined as

$$
W\left(S_{N}\right)=\max _{a, b, t}\left|\sum_{j=0}^{t-1} s_{a+j b}\right|
$$

where the maximum is taken over all $a, b, t$ such that $a, b, t \in \mathbf{N}$ and $1 \leq a \leq a+(t-1) b \leq N$, while the correlation measure of order $k$ (or order $k$ correlation measure) of $S_{N}$ is defined as

$$
C_{k}\left(S_{N}\right)=\max _{M, D}\left|\sum_{n=1}^{M} s_{n+d_{1}} s_{n+d_{2}} \cdots s_{n+d_{k}}\right|
$$

[^0]where the maximum is taken over all $D=\left(d_{1}, \cdots, d_{k}\right)$ with non-negative integers $0 \leq d_{1}<$ $\cdots<d_{k}$ and $M$ such that $M+d_{k} \leq N$.
$S_{N}$ is considered as a "good" pseudorandom sequence, if both $W\left(S_{N}\right)$ and $C_{k}\left(S_{N}\right)$ (at least for small $k$ ) are "small" in terms of $N$ (in particular, both are $o(N)$ as $N \rightarrow \infty$ ). It was shown in [2] that for a "truly" random sequence $S_{N} \in\{+1,-1\}^{N}$ (i.e., choosing $S_{N} \in\{+1,-1\}^{N}$ with probability $1 / 2^{N}$ ), both $W\left(S_{N}\right)$ and $C_{k}\left(S_{N}\right)$ (for some fixed $k$ ) are around $N^{1 / 2} \log ^{c}(N)$ with "near 1" probability. From [1] the Legendre sequence forms a "good" pseudorandom sequence. Many other "good" (but slightly inferior) binary sequences were designed in the literature, see for example [1-6] and references therein.

Very recently, Liu ${ }^{[7,8]}$, Louboutin, Rivat and Sárközy ${ }^{[9]}$ present some constructions of finite binary sequences (see below) related to Lehmer numbers respectively. These sequences are shown to be "good" pseudorandom sequences.

Let $p$ be an odd prime. We denote by $F_{p}=\{0,1, \cdots, p-1\}$ the finite field of $p$ elements, by $F_{p}^{*}$ the multiplicative group of $F_{p}$. Let

$$
\bar{\gamma}= \begin{cases}\gamma^{-1}, & \text { if } \gamma \in F_{p}^{*} \\ 0, & \text { if } \gamma=0\end{cases}
$$

For any $\gamma \in F_{p}^{*}$, we always suppose that $\bar{\gamma} \in\{1, \cdots, p-1\}$.
Construction 1. ${ }^{[9]}$ Let $f(x) \in F_{p}[x]$ be of degree $d$ with $1 \leq d<p$. Suppose $f(n), \overline{f(n)} \in$ $\{0, \cdots, p-1\}$ for any $n \in F_{p}$. Define the sequence $E_{p-1}=\left\{s_{1}, s_{2}, \cdots, s_{p-1}\right\}$ by

$$
s_{n}:=\left\{\begin{array}{lll}
+1, & \text { if } f(n) \equiv \overline{f(n)} & (\bmod 2) \\
-1, & \text { if } f(n) \not \equiv \overline{f(n)} & (\bmod 2)
\end{array}\right.
$$

Then the bounds hold

$$
W\left(E_{p-1}\right) \ll(d+s) p^{1 / 2} \log ^{3}(p) \text { and } C_{2}\left(E_{p-1}\right) \ll(d+s) p^{1 / 2} \log ^{5}(p)
$$

where $s$ is the number of distinct roots of $f(x)$ in an algebraic closure of $F_{p}$.
It is easy to see for $n \in\{1, \cdots, p-1\}, s_{n}=(-1)^{f(n)+\overline{f(n)}}$.
Construction 2. ${ }^{[8]}$ Suppose $\bar{\gamma} \in\{0, \cdots, p-1\}, \forall \gamma \in F_{p}$. Define the sequence $E_{p-1}^{\prime}=$ $\left\{s_{1}^{\prime}, s_{2}^{\prime}, \cdots, s_{p-1}^{\prime}\right\}$ by

$$
s_{n}^{\prime}:= \begin{cases}(-1)^{n+\bar{n}}, & \text { if } n \text { is a quadratic residue } \bmod p \\ (-1)^{n+\bar{n}+1}, & \text { if } n \text { is a quadratic nonresidue } \bmod p\end{cases}
$$

Then the bounds hold $W\left(E_{p-1}^{\prime}\right) \ll p^{1 / 2} \log ^{2}(p)$ and $C_{2}\left(E_{p-1}^{\prime}\right) \ll p^{1 / 2} \log ^{3}(p)$.
Construction 3. ${ }^{[7]}$ Suppose $\bar{\gamma} \in\{0, \cdots, p-1\}$ for any $\gamma \in F_{p}$. Let $c \in\{1, \cdots, p-1\}$ be a fixed number. Define the sequence $E_{p-1}^{\prime \prime}=\left\{s_{1}^{\prime \prime}, s_{2}^{\prime \prime}, \cdots, s_{p-1}^{\prime \prime}\right\}$ by

$$
s_{n}^{\prime \prime}:= \begin{cases}(-1)^{\bar{n}+\overline{n+c}}, & \text { if } p \nmid n(n+c) \\ 1, & \text { otherwise }\end{cases}
$$

Then the bounds hold $W\left(E_{p-1}^{\prime \prime}\right) \ll p^{1 / 2} \log ^{3}(p)$ and $C_{2}\left(E_{p-1}^{\prime \prime}\right) \ll p^{1 / 2} \log ^{5}(p)$.

We note that in the above three constructions, the numbers

$$
f(n)+\overline{f(n)}, n+\bar{n}, \bar{n}+\overline{n+c} \in\{0, \cdots, 2 p-2\}
$$

for all $n \in\{1, \cdots, p-1\}$. In particular, it seems difficult to consider the correlation measure of higher order ( $>2$ ).

Motivated by these constructions, we present two new constructions. One is a variation of Construction 3, the other is constructed using power functions. We estimate their welldistribution measure and correlation measure of order $k(\geq 2)$, which indicate that the resulting sequences may form "good" pseudorandom sequences.

Throughout this paper, the implied constants in the symbol " $\ll$ " are absolute.

## §2 New constructions

In our constructions, we will use 0 and 1 to represent the terms of a binary sequence.
Definition 1. Suppose $\bar{\gamma} \in\{0, \cdots, p-1\}$ for any $\gamma \in F_{p}$. Let $c \in\{1, \cdots, p-1\}$ be a fixed number. Define the sequence $\mathcal{X}_{p-1}=\left\{x_{1}, x_{2}, \cdots, x_{p-1}\right\}$ by

$$
x_{n}:= \begin{cases}1, & \text { if }[\bar{n}+\overline{n+c}]_{p} \text { is odd } \\ 0, & \text { if }[\bar{n}+\overline{n+c}]_{p} \text { is even }\end{cases}
$$

where $[u]_{p}$ denotes the unique $r \in\{0,1, \cdots, p-1\}$ such that $u \equiv r \quad(\bmod p)$.
Definition 2. Suppose $\bar{\gamma} \in\{0, \cdots, p-1\}$ for any $\gamma \in F_{p}$ and $c \in\{1, \cdots, p-1\}$ is a fixed number. Let $\eta \in F_{p}^{*}$ be a fixed primitive element. Define the sequence $\mathcal{Y}_{p-1}=\left\{y_{1}, y_{2}, \cdots, y_{p-1}\right\}$ by

$$
y_{n}:=\left\{\begin{array}{lll}
0, & \text { if } \overline{\eta^{n}} \equiv \overline{\eta^{n}+c} & (\bmod 2), \\
1, & \text { if } \overline{\eta^{n}} \not \equiv \overline{\eta^{n}+c} & (\bmod 2) .
\end{array}\right.
$$

### 2.1. Well-distribution and correlation of $\mathcal{X}_{p-1}$

Theorem 1. Let $\mathcal{X}_{p-1}=\left\{x_{1}, x_{2}, \cdots, x_{p-1}\right\}$ be defined as in Definition 1. Then we have

$$
W\left(\mathcal{X}_{p-1}\right) \ll p^{1 / 2} \log ^{2}(p)
$$

and

$$
C_{k}\left(\mathcal{X}_{p-1}\right) \ll 2^{k} k p^{1 / 2} \log ^{k+1}(p)
$$

for $k<p$. In particular, $C_{2}\left(\mathcal{X}_{p-1}\right) \ll p^{1 / 2} \log ^{3}(p)$.
In order to prove Theorem 1, we need some basic results. Let $e_{m}(z)=\exp (2 \pi i z / m)$.
Lemma 1. ${ }^{[10,11]}$ Let $m>1$ be a positive integer. Then

$$
\sum_{c=0}^{m-1}\left|\sum_{z=H+1}^{H+N} e_{m}(c z)\right| \leq m(1+\log (m))
$$

holds for any integers $H$ and $1 \leq N \leq m$.

Lemma 2. Let $p$ be an odd prime number and $\lambda \in \mathbf{Z}$ with $0 \leq|\lambda| \leq \frac{p-1}{2}$. We define

$$
\begin{equation*}
U(\lambda):=\sum_{r=0}^{(p-1) / 2} e_{p}(-2 \lambda r)-\sum_{r=1}^{(p-1) / 2} e_{p}(2 \lambda r) . \tag{1}
\end{equation*}
$$

Then the following bound holds: $\sum_{|\lambda| \leq(p-1) / 2}|U(\lambda)| \leq 2 p(1+\log (p))$.
Proof. Since $|U(\lambda)| \leq\left|\sum_{r=0}^{(p-1) / 2} e_{p}(-2 \lambda r)\right|+\left|\sum_{r=1}^{(p-1) / 2} e_{p}(2 \lambda r)\right|$, the desired result follows from Lemma 1.

Lemma 3. ${ }^{[7,12]}$ For any polynomials $g(x), h(x) \in F_{p}[x]$ such that the rational function $F(x)=$ $g(x) / h(x)$ is not constant on $F_{p}$, let $\chi$ be a nontrivial multiplicative character of $F_{p}$ and $s$ the number of distinct roots of the polynomial $h(x)$ in an algebraic closure of $F_{p}$. For $p \nmid \lambda$, we have
(i)
(i) $\left|\sum_{\substack{\xi \in F_{p} \\ h(\xi) \neq 0}} e_{p}(\lambda F(\xi))\right| \leq\left(\max (\operatorname{deg}(g), \operatorname{deg}(h))+s^{*}-2\right) \sqrt{p}+\delta$, where $s^{*}=s$ and $\delta=1$ if $\operatorname{deg}(g) \leq \operatorname{deg}(h)$, and $s^{*}=s+1$ and $\delta=0$ otherwise.
(ii)

$$
\left\lvert\, \begin{aligned}
& \sum_{\substack{\xi \in F^{*} \\
h(\xi) \neq 0}} e_{p}(\lambda F(\xi)) \chi(\xi) \mid \leq\left(\max (\operatorname{deg}(g), \operatorname{deg}(h))+s^{*}-1\right) \sqrt{p} \text {, where } s^{*}=s \text { if } \operatorname{deg}(g) \leq \\
& \operatorname{deg}(h) \text {, and } s^{*}=s+1 \text { otherwise. }
\end{aligned}\right.
$$

Remark 1. Let $F(x)$ be defined as in Lemma 3. According to Lemma 3 and Lemma 1, one can estimate incomplete sums

$$
\sum_{\substack{n=A+1 \\ h(n) \neq 0}}^{B} e_{p}(F(n)) \text { and } \sum_{\substack{n=A+1 \\ h\left(\eta^{n}\right) \neq 0}}^{B} e_{p}\left(F\left(\eta^{n}\right)\right),
$$

where $0 \leq A<B \leq p-1$ and $\eta \in F_{p}^{*}$ is a primitive element. In fact,

$$
\begin{aligned}
& \sum_{\substack{n=A+1 \\
h(n \neq 0}}^{B} e_{p}(F(n))=\sum_{n=A+1}^{B} \sum_{\substack{m=0 \\
h-1}}^{p-1} e_{p}(F(m)) \cdot \frac{1}{p} \sum_{\mu=0}^{p-1} e_{p}(\mu(n-m)) \\
= & \frac{1}{p} \sum_{\mu=0} \sum_{n=A+1}^{B} e_{p}(\mu n) \sum_{\substack{m=0 \\
h(m) \neq 0}}^{p-1} e_{p}(F(m)-\mu m),
\end{aligned}
$$

and

$$
\begin{aligned}
& \sum_{\substack{n=A+1 \\
h\left(\eta^{n} \neq 0\right.}}^{B} e_{p}\left(F\left(\eta^{n}\right)\right)=\sum_{n=A+1}^{B} \sum_{\substack{m=1 \\
n=1 \\
p-1}}^{p-1} e_{p}\left(F\left(\eta^{m}\right)\right) \cdot \frac{1}{p-1} \sum_{\mu=1}^{p-1} e_{p-1}(\mu(m-n)) \\
= & \frac{1}{p-1} \sum_{\mu=1}^{p-1} \sum_{n=A+1}^{B} e_{p-1}(-\mu n) \sum_{\substack{\xi=1 \\
h(\xi) \neq 0}} e_{p}(F(\xi)) \chi_{\mu}(\xi),
\end{aligned}
$$

where $\chi_{\mu}\left(\eta^{m}\right)=e_{p-1}(\mu m)$ is a multiplicative character of $F_{p}$.

Proof of Theorem 1. Let $f(n)=\bar{n}+\overline{n+c}$. For any $1 \leq n \leq p-1$ with $p \nmid(n+c)$ we have

$$
\frac{1}{p} \sum_{r=0}^{(p-1) / 2} \sum_{|\lambda| \leq(p-1) / 2} e_{p}(\lambda(f(n)-2 r))= \begin{cases}1, & \text { if }[f(n)]_{p} \text { is even }  \tag{2}\\ 0, & \text { if }[f(n)]_{p} \text { is odd }\end{cases}
$$

and

$$
\frac{1}{p} \sum_{r=1}^{(p-1) / 2} \sum_{|\lambda| \leq(p-1) / 2} e_{p}(\lambda(f(n)+2 r))= \begin{cases}0, & \text { if }[f(n)]_{p} \text { is even }  \tag{3}\\ 1, & \text { if }[f(n)]_{p} \text { is odd }\end{cases}
$$

Subtracting (3) from (2) yields

$$
\frac{1}{p} \sum_{|\lambda| \leq(p-1) / 2} e_{p}(\lambda f(n)) U(\lambda)= \begin{cases}1, & \text { if }[f(n)]_{p} \text { is even } \\ -1, & \text { if }[f(n)]_{p} \text { is odd }\end{cases}
$$

where $U(\lambda)$ is defined as in Lemma 2.
It is easy to see that for any $1 \leq n \leq p-1$ with $p \nmid(n+c)$,

$$
\begin{equation*}
(-1)^{x_{n}}=\frac{1}{p} \sum_{|\lambda| \leq(p-1) / 2} e_{p}(\lambda f(n)) U(\lambda) . \tag{4}
\end{equation*}
$$

For $a, b, t \in \mathbf{N}$ with $1 \leq a \leq a+(t-1) b \leq p-1$, we have

$$
\left.\begin{aligned}
& \left|\sum_{j=0}^{t-1}(-1)^{x_{a+j}}\right| \leq\left|\sum_{\substack{j=0 \\
p \nmid a+j b+c}}^{t-1}(-1)^{x_{a+j b}}\right|+1=\frac{1}{p}\left|\sum_{\substack{j=0 \\
p \nmid a+j b+c}}^{t-1} \sum_{|\lambda| \leq(p-1) / 2} U(\lambda) e_{p}(\lambda f(a+j b))\right|+1 \\
= & \frac{1}{p}\left|\sum_{|\lambda| \leq(p-1) / 2} U(\lambda) \sum_{\substack{j=0 \\
p \nmid a+j b+c}}^{t-1} e_{p}(\lambda f(a+j b))\right|+1 \\
\leq & \frac{1}{p} \sum_{|\lambda| \leq(p-1) / 2}|U(\lambda)| \cdot\left|\sum_{\substack{j=0 \\
p \nmid a+j b+c}}^{t-1} e_{p}(\lambda f(a+j b))\right|+1 \\
\leq & \frac{1}{p}\left(\sum_{|\lambda|=1}^{(p-1) / 2}|U(\lambda)| \cdot\left|\sum_{\substack{j=0 \\
p \nmid a+j b+c}}^{t-1} e_{p}(\lambda f(a+j b))\right|+t\right)+1 \\
\leq & \left.\frac{1}{p} \sum_{|\lambda|=1}^{(p-1) / 2}|U(\lambda)| \cdot \right\rvert\, \sum_{\substack{j=0 \\
p-1}}^{t a+j b+c} ⿺
\end{aligned} e_{p}(\lambda f(a+j b)) \right\rvert\,+2 .
$$

Now by Lemmas 2 and 3, we obtain the bound of $W\left(\mathcal{X}_{p-1}\right)$.
For $D=\left(d_{1}, \cdots, d_{k}\right)$ and $M$ with $0 \leq d_{1}<\cdots<d_{k} \leq p-1-M$, there are at most $k$ elements $n(1 \leq n \leq M<p-1)$ such that $p \mid\left(n+d_{j}+c\right)$ for at least one number $j: 1 \leq j \leq k$.

Then we have

$$
\begin{aligned}
& \left|\sum_{n=1}^{M}(-1)^{x_{n+d_{1}}+\cdots+x_{n+d_{k}}}\right| \leq\left|\sum_{\substack{\left.n=1 \\
p \nmid n+d_{j}+x\right) \\
1 \leq j \leq k}}^{M} \prod_{i=1}^{k}\left(\frac{1}{p} \sum_{\left|\lambda_{i}\right| \leq(p-1) / 2} U\left(\lambda_{i}\right) e_{p}\left(\lambda_{i} f\left(n+d_{i}\right)\right)\right)\right|+k \\
= & \frac{1}{p^{k}}\left|\sum_{\left|\lambda_{1}\right| \leq(p-1) / 2} \cdots \sum_{\left|\lambda_{k}\right| \leq(p-1) / 2} U\left(\lambda_{1}\right) \cdots U\left(\lambda_{k}\right) \sum_{\substack{n=1 \\
p \nmid\left(n+d_{j}+x\right) \\
1 \leq j \leq k}}^{M} e_{p}\left(\sum_{i=1}^{k} \lambda_{i} f\left(n+d_{i}\right)\right)\right|+k \\
\leq & \frac{1}{p^{k}}\left(\left(\sum_{\substack{ \\
0<\left|\lambda_{i}\right| \leq(p-1) / 2}}\left|U\left(\lambda_{i}\right)\right|\right)^{k} \cdot\left|\sum_{\substack{n=1 \\
p \nmid\left(n+d_{j}+c\right) \\
1 \leq j \leq k}}^{M} e_{p}\left(\sum_{i=1}^{k} \lambda_{i} f\left(n+d_{i}\right)\right)\right|+M\right)+k .
\end{aligned}
$$

If $\sum_{i=1}^{k} \lambda_{i} f\left(y+d_{i}\right)$ is a nonconstant rational function in $F_{p}(y)$ when $\lambda_{1}, \cdots, \lambda_{k}$ are not all zero, then by Lemma 3, $\left|\sum_{\substack{\left.n=1 \\ p \nmid n+d_{j}+c\right) \\ 1 \leq j \leq k}}^{M} e_{p}\left(\sum_{i=1}^{k} \lambda_{i} f\left(n+d_{i}\right)\right)\right| \leq 4 k \sqrt{p}(1+\log (p))$. So

$$
\left|\sum_{n=1}^{M}(-1)^{x_{n+d_{1}}+\cdots+x_{n+d_{k}}}\right| \leq 4 k 2^{k} \sqrt{p}(1+\log (p))^{k+1}+k+1=O\left(k 2^{k} \sqrt{p} \log ^{k+1}(p)\right)
$$

It remains to prove if $\lambda_{1}, \cdots, \lambda_{k}$ are not all zero, $\sum_{i=1}^{k} \lambda_{i} f\left(y+d_{i}\right)$ is a nonconstant rational function. Suppose that there are $s(1 \leq s \leq k<p)$ elements $\lambda_{i 1}, \cdots, \lambda_{i s}$ are not zero (while other coefficients are zero), i.e.,

$$
F(y) \triangleq \sum_{i=1}^{k} \lambda_{i} f\left(y+d_{i}\right)=\lambda_{i 1} f\left(y+d_{i 1}\right)+\lambda_{i 2} f\left(y+d_{i 2}\right)+\cdots+\lambda_{i s} f\left(y+d_{i s}\right)
$$

Let

$$
H(y)=\left(y+d_{i 1}\right)\left(y+d_{i 1}+c\right)\left(y+d_{i 2}\right)\left(y+d_{i 2}+c\right) \cdots\left(y+d_{i s}\right)\left(y+d_{i s}+c\right) \in F_{p}[y] .
$$

If $\beta \in F_{p}$ is a zero of $H(y)$ and $(y-\beta)^{2} \nmid H(y)$, then $\beta$ is a pole of $F(y)$, therefore $F(y)$ is nonconstant. While if

$$
H(y)=\left(y+d_{i 1}\right)^{2}\left(y+d_{i 2}\right)^{2} \cdots\left(y+d_{i s}\right)^{2}
$$

then since $d_{i 1}<d_{i 2}<\cdots<d_{i s}$, we have

$$
\begin{cases}d_{i 1} \equiv d_{j 1}+c & (\bmod p) \\ d_{i 2} \equiv d_{j 2}+c & (\bmod p) \\ \cdots \\ d_{i s} \equiv d_{j s}+c & (\bmod p)\end{cases}
$$

where $d_{j 1}, d_{j 2}, \cdots, d_{j s}$ is a permutation of $d_{i 1}, d_{i 2}, \cdots, d_{i s}$. So $s c \equiv 0 \quad(\bmod p)$, which yields $c=0$. it is a contradiction since $c \in\{1, \cdots, p-1\}$. Therefore, there exists a zero $\beta$ of $H(y)$ such that $(y-\beta)^{2} \nmid H(y)$, which makes $F(y)$ to be nonconstant. The proof is completed.

### 2.2. Well-distribution and correlation of $\mathcal{Y}_{p-1}$

Theorem 2. Let $\mathcal{Y}_{p-1}=\left\{y_{1}, y_{2}, \cdots, y_{p-1}\right\}$ be defined as in Definition 2. Then we have

$$
W\left(\mathcal{Y}_{p-1}\right) \ll p^{1 / 2} \log ^{3}(p) \quad \text { and } \quad C_{2}\left(\mathcal{Y}_{p-1}\right) \ll p^{1 / 2} \log ^{5}(p)
$$

Lemma 4. $\sum_{|\lambda| \leq(p-1) / 2}\left|\sum_{u=1}^{p-1}(-1)^{u} e_{p}(-\lambda u)\right| \leq 2 p(1+\log (p))$.
Proof. In fact,

$$
\sum_{u=1}^{p-1}(-1)^{u} e_{p}(-\lambda u)=\sum_{r=1}^{(p-1) / 2} e_{p}(-2 \lambda r)-\sum_{r=1}^{(p-1) / 2} e_{p}(-\lambda(2 r-1))
$$

Then the proof is similar to that of Lemma 2.
Proof of Theorem 2. Our approach follows the path of $[7,8]$. It is easy to see for $n \in$ $\{1, \cdots, p-1\}$ with $p \nmid\left(\eta^{n}+c\right)$,

$$
\begin{equation*}
(-1)^{y_{n}}=(-1)^{\overline{\eta^{n}}+\overline{\eta^{n}+c}} \tag{5}
\end{equation*}
$$

For $a, b, t \in \mathbf{N}$ with $1 \leq a \leq a+(t-1) b \leq p-1$, we have

$$
\begin{aligned}
& \quad\left|\sum_{j=0}^{t-1}(-1)^{y_{a+j b}}\right| \leq\left|\sum_{\substack{j=0 \\
p \nmid \eta^{a+j b}+c}}^{t-1}(-1)^{\overline{\eta^{a+j b}}+\overline{\eta^{a+j b}+c}}\right|+1 \\
& \left.=\frac{1}{p^{2}} \right\rvert\, \sum_{\substack{j=0 \\
p \nmid \eta^{+j b}+c}}^{t-1} \sum_{u=1}^{p-1} \sum_{|\lambda| \leq(p-1) / 2} e_{p}\left(\lambda\left(\overline{\eta^{a+j b}}-u\right)\right) \\
& \quad \times \frac{1}{p-1} \sum_{v=1}^{p^{2}}\left|\sum_{|\lambda| \leq(p-1) / 2} e_{p}\left(\mu\left(\overline{\eta^{a+j b}+c}-v\right)\right)(-1)^{u+v}\right|+1 \\
& \quad \sum_{|\lambda| \leq(p-1) / 2}^{p-1}(-1)^{u} e_{p}(-\lambda u) \\
& \quad \times \sum_{|\mu| \leq(p-1) / 2} \sum_{v=1}^{p-1}(-1)^{v} e_{p}(-\mu v) \sum_{\substack{j=0 \\
p \nmid \eta^{a+j b}+c}}^{t-1} e_{p}\left(\lambda \overline{\eta^{a+j b}}+\mu \overline{\eta^{a+j b}+c}\right) \mid+1
\end{aligned}
$$

Suppose the multiplicative order of $\eta^{b} \in F_{p}^{*}$ is $T$. For $\lambda \neq 0$ and $\mu \neq 0$, by Lemma 3 we have

$$
\begin{aligned}
& \left|\sum_{\substack{j=0 \\
p \nmid \eta^{a+j b}+c}}^{T-1} e_{p}\left(\lambda \overline{\eta^{a+j b}}+\mu \overline{\eta^{a+j b}+c}\right)\right|=\left|\sum_{\substack{j=0 \\
p \nmid \eta^{a+j b}+c}}^{T-1} e_{p}\left(\lambda \overline{\eta^{a}\left(\eta^{b}\right)^{j}}+\mu \overline{\eta^{a}\left(\eta^{b}\right)^{j}+c}\right)\right| \\
= & \frac{T}{p-1}\left|\sum_{\xi \in F_{p}^{*}}^{*} e_{p}\left(\frac{\lambda\left(\eta^{a} \xi^{(p-1) / T}+c\right)+\mu \eta^{a} \xi^{(p-1) / T}}{\eta^{a} \xi^{(p-1) / T}\left(\eta^{a} \xi^{(p-1) / T}+c\right)}\right)\right| \leq 4 p^{1 / 2}+1,
\end{aligned}
$$

where $\sum^{*}$ indicates that the poles of the corresponding rational function are excluded from the summation. So by Remark 1 we have

$$
\left|\sum_{\substack{j=0 \\ p \nmid \eta^{a+j b}+c}}^{t-1} e_{p}\left(\lambda \overline{\eta^{a+j b}}+\mu \overline{\eta^{a+j b}+c}\right)\right| \leq(4 \sqrt{p}+1)(1+\log (p))
$$

Now by Lemma 4 we have

$$
\left|\sum_{j=0}^{t-1}(-1)^{y_{a+j b}}\right| \leq 4(4 \sqrt{p}+1)(1+\log (p))^{3}
$$

For integers $d_{1}, d_{2}$ and $M$ with $0 \leq d_{1}<d_{2} \leq p-1-M$, we have

$$
\begin{aligned}
& \left|\sum_{n=1}^{M}(-1)^{y_{n+d_{1}}+y_{n+d_{2}}}\right| \leq\left|\sum_{\substack{n=1 \\
p \nmid\left(\eta^{n+d_{1}}+c\right)\left(\eta^{n+d_{2}}+c\right)}}^{M}(-1)^{\overline{\eta^{n+d_{1}}}+\overline{\eta^{n+d_{1}+c}+} \overline{\eta^{n+d_{2}}}+\overline{\eta^{n+d_{2}}+c}}\right|+2 \\
& \left.=\frac{1}{p^{4}} \right\rvert\, \sum_{\substack{n=1 \\
p \nmid\left(\eta^{n+d_{1}}+c\right)\left(\eta^{n+d_{2}}+c\right)}}^{M} \sum_{u_{1}=1}^{p-1} \sum_{\left|\lambda_{1}\right| \leq(p-1) / 2} e_{p}\left(\lambda_{1}\left(\overline{\eta^{n+d_{1}}}-u_{1}\right)\right) \\
& \times \sum_{u_{2}=1}^{p-1} \sum_{\left|\lambda_{2}\right| \leq(p-1) / 2} e_{p}\left(\lambda_{2}\left(\overline{\eta^{n+d_{1}}+c}-u_{2}\right)\right) \sum_{u_{3}=1}^{p-1} \sum_{\left|\lambda_{3}\right| \leq(p-1) / 2} e_{p}\left(\lambda_{3}\left(\overline{\eta^{n+d_{2}}}-u_{3}\right)\right) \\
& \times \sum_{u_{4}=1}^{p-1} \sum_{\left|\lambda_{4}\right| \leq(p-1) / 2} e_{p}\left(\lambda_{4}\left(\overline{\eta^{n+d_{2}}+c}-u_{4}\right)\right) \cdot(-1)^{u_{1}+u_{2}+u_{3}+u_{4}} \mid+2 \\
& \left.=\frac{1}{p^{4}} \right\rvert\, \sum_{\left|\lambda_{1}\right| \leq(p-1) / 2} \sum_{u_{1}=1}^{p-1}(-1)^{u_{1}} e_{p}\left(-\lambda_{1} u_{1}\right) \cdot \sum_{\left|\lambda_{2}\right| \leq(p-1) / 2} \sum_{u_{2}=1}^{p-1}(-1)^{u_{2}} e_{p}\left(-\lambda_{2} u_{2}\right) \\
& \times \sum_{\left|\lambda_{3}\right| \leq(p-1) / 2} \sum_{u_{3}=1}^{p-1}(-1)^{u_{3}} e_{p}\left(-\lambda_{3} u_{3}\right) \cdot \sum_{\left|\lambda_{4}\right| \leq(p-1) / 2} \sum_{u_{4}=1}^{p-1}(-1)^{u_{4}} e_{p}\left(-\lambda_{4} u_{4}\right) \\
& \times \sum_{\substack{n=1 \\
p \nmid\left(\eta^{n+d_{1}}+c\right)\left(\eta^{n+d_{2}}+c\right)}}^{M} e_{p}\left(\lambda_{1} \overline{\eta^{n+d_{1}}}+\lambda_{2} \overline{\eta^{n+d_{1}}+c}+\lambda_{3} \overline{\eta^{n+d_{2}}}+\lambda_{4} \overline{\eta^{n+d_{2}}+c}\right) \mid+2 .
\end{aligned}
$$

Since $\eta$ is a primitive element of $F_{p}^{*}$, for $\lambda_{1}, \cdots, \lambda_{4}$ are not all zero, by Lemma 3 we have

$$
\begin{aligned}
& \quad \sum_{n=1}^{p-1} e_{p}\left(\lambda_{1} \overline{\eta^{n+d_{1}}}+\lambda_{2} \overline{\eta^{n+d_{1}}+c}+\lambda_{3} \overline{\eta^{n+d_{2}}}+\lambda_{4} \overline{\eta^{n+d_{2}}+c}\right) \\
& p \nmid\left(\eta^{\left.n+d_{1}+c\right)\left(\eta^{n+d_{2}}+c\right)}\right. \\
& =\sum_{\xi \in F_{p}^{*}}^{*} e_{p}\left(\lambda_{1} \overline{\eta^{d_{1}} \xi}+\lambda_{2} \overline{\eta^{d_{1}} \xi+c}+\lambda_{3} \overline{\eta^{d_{2}} \xi}+\lambda_{4} \overline{\eta^{d_{2}} \xi+c}\right) \leq 5 \sqrt{p}+1
\end{aligned}
$$

where $\sum^{*}$ indicates that the poles of the corresponding rational function are excluded from the summation. Now by Remark 1 we have

$$
\begin{aligned}
& \sum_{\substack{n=1 \\
M}}^{M} e_{p}\left(\lambda_{1} \overline{\eta^{n+d_{1}}}+\lambda_{2} \overline{\eta^{n+d_{1}}+c}+\lambda_{3} \overline{\eta^{n+d_{2}}}+\lambda_{4} \overline{\eta^{n+d_{2}}+c}\right) \mid \\
& \leq(5 \sqrt{p}+1)(1+\log (p)) .
\end{aligned}
$$

So by Lemma 4 we derive

$$
\left|\sum_{n=1}^{M}(-1)^{y_{n+d_{1}}+y_{n+d_{2}}}\right| \leq 8(5 \sqrt{p}+1)(1+\log (p))^{5}+2
$$

which completes the proof.

### 2.3. Linear complexity profile

We recall that the linear complexity profile of a binary sequence

$$
S=\left\{s_{0}, s_{1}, \cdots\right\} \in\{0,1\}^{\infty}
$$

is the function $L(S, N)$ defined for every positive integer $N$, as the least order $l$ of a linear recurrence relation

$$
s_{n}=c_{1} s_{n-1}+\cdots+c_{l} s_{n-l}, \quad c_{i}=0,1
$$

for all $n$ with $l \leq n \leq N-1$, which $S$ satisfies. We use the convention that $L(S, N)=0$ if the first $N$ elements of $S$ are all zero and $L(S, N)=N$ if the first $N-1$ elements of $S$ are zero and $s_{N-1}=1$. The value

$$
L(S)=\sup _{N \geq 1} L(S, N)
$$

is called the linear complexity of $S$, see for example [13]. For the linear complexity of any periodic sequence of period $t$ one easily verifies that $L(S)=L(S, 2 t) \leq t$. It is desirable to have sequences with large linear complexity for cryptographic applications.

Proposition 1. ${ }^{[14]}$ Let $S$ be a $T$-periodic binary sequence. For $2 \leq N \leq T-1$ we have

$$
L(S, N) \geq N-\max _{1 \leq k \leq L(S, N)+1} C_{k}(S)
$$

Corollary 1. Let $\mathcal{X}_{p-1}=\left\{x_{1}, x_{2}, \cdots, x_{p-1}\right\}$ be defined as in Definition 1. For $2 \leq N \leq p-1$ we have

$$
L\left(\mathcal{X}_{p-1}, N\right)=\Omega\left(\frac{\log \left(N / p^{3 / 4}\right)}{\log \log (p)}\right)
$$

Proof. The proof is similar to that of [14,Corollary 1], we give below for completeness. From Proposition 1 (see the proof of $[14$,Theorem 1]), we see that

$$
N-L\left(\mathcal{X}_{p-1}, N\right) \leq \max _{1 \leq k \leq L\left(\mathcal{X}_{p-1}, N\right)+1} C_{k}\left(\mathcal{X}_{p-1}\right)
$$

which yields

$$
N \ll L 2^{L} \sqrt{p} \log ^{L+2}(p)
$$

by Theorem 1. Suppose $L \leq \log \left(p^{1 / 4}\right)$, otherwise the result is trivial. Then we have

$$
N \ll p^{3 / 4} \log ^{L+3}(p)
$$

hence we obtain

$$
L \gg \frac{\log \left(N / p^{3 / 4}\right)}{\log \log (p)}
$$

Table 1 Comparison of our sequences with some other sequences

| Sequences | Length | Well-distributied | Correlation of order $k$ |
| :--- | :--- | :--- | :--- |
| Legendre sequence | $N=p-1$ | $O\left(p^{1 / 2} \log (p)\right)$ | $O\left(k p^{1 / 2} \log (p)\right) ; k \geq 2$ |
| index sequence | $N=p-1$ | $O\left(p^{1 / 2} \log ^{2}(p)\right)$ | $O\left(k 4^{k} p^{1 / 2} \log ^{k+1}(p)\right) ; k \geq 2$ |
| $E_{p-1}$ | $N=p-1$ | $O\left(p^{1 / 2} \log ^{3}(p)\right)$ | $O\left(p^{1 / 2} \log ^{5}(p)\right) ; k=2$ |
| $E_{p-1}^{\prime \prime}$ | $N=p-1$ | $O\left(p^{1 / 2} \log ^{2}(p)\right)$ | $O\left(p^{1 / 2} \log ^{3}(p)\right) ; k=2$ |
| $E_{p-1}^{\prime \prime}$ | $N=p-1$ | $O\left(p^{1 / 2} \log ^{3}(p)\right)$ | $O\left(p^{1 / 2} \log ^{5}(p)\right) ; k=2$ |
| $\mathcal{X}_{p-1}$ | $N=p-1$ | $O\left(p^{1 / 2} \log ^{2}(p)\right)$ | $O\left(k 2^{k} p^{1 / 2} \log ^{k+1}(p)\right) ; k \geq 2$ |
| $\mathcal{Y}_{p-1}$ | $N=p-1$ | $O\left(p^{1 / 2} \log ^{3}(p)\right)$ | $O\left(p^{1 / 2} \log ^{5}(p)\right) ; k=2$ |

Remark 2. The implied constant in the symbol " $O$ " may sometimes depend on the degree $\operatorname{deg}(f)$ of a function $f$ adopted in the corresponding constructions and is absolute otherwise.

## §3 Conclusions

We have constructed two families of finite binary sequences using multiplicative inverse, which were used in $[7,8,9]$ to construct different sequences described in Constructions 1,2 and 3 , respectively. The sequence $\mathcal{X}_{p-1}$ is a variation of Construction 3, while the sequence $\mathcal{Y}_{p-1}$ is constructed using power functions. Two important pseudorandom measures, the well-distribution measure and the correlation measure of order $k$, are estimated by using some estimates of certain exponential sums.

In Table 1, we compare our sequences with some other sequences, such as the Legendre sequence ${ }^{[1]}$, the index sequence ${ }^{[4]}$ and $E_{p-1}, E_{p-1}^{\prime}, E_{p-1}^{\prime \prime}$ described in $\S 1$. We conclude that our sequences also have strong pseudo-random properties. So these constructions may provide a very attractive alternative to traditional methods in applications.

From a point of implementation view, the sequences can be computed fast, since the multiplicative inverse can be computed fast (in polynomial time).

Finally we remark that in $[15,16]$ the recursive inversive generators, explicit inversive generators and explicit nonlinear generators are used to build families of binary sequences with strong pseudorandom properties in a different way.

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