# Algebraic Properties and Panconnectivity of Folded Hypercubes<sup>\*</sup>

Meijie Ma $^{a\dagger}$  Jun-Ming Xu $^{b}$ 

<sup>a</sup>School of Mathematics and System Science, Shandong University Jinan, 250100, China
<sup>b</sup>Department of Mathematics, University of Science and Technology of China Hefei, 230026, China

**Abstract** This paper considers the folded hypercube  $FQ_n$ , as an enhancement on the hypercube, and obtain some algebraic properties of  $FQ_n$ . Using these properties the authors show that for any two vertices x and y in  $FQ_n$  with distance d and any integers  $h \in \{d, n + 1 - d\}$  and l with  $h \leq l \leq 2^n - 1$ ,  $FQ_n$  contains an xy-path of length l and no xy-paths of other length provided that l and h have the same parity.

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### 1 Introduction

It is well-known that a topological structure for an interconnection network can be modelled by a connected graph G = (V, E) [14]. As a topology for an interconnection network of a multiprocessor system, the hypercube structure is a widely used and well-known interconnection model since it possesses many attractive properties [8, 14]. The *n*-dimensional hypercube  $Q_n$  is a graph with  $2^n$  vertices, each vertex with a distinct binary string  $x_1x_2\cdots x_n$  of length *n* on the set  $\{0, 1\}$ , and two vertices being linked by an edge if and only if their strings differ in exactly one bit.

As a variant of the hypercube, the *n*-dimensional folded hypercube  $FQ_n$ , proposed first by El-Amawy and Latifi [3], is a graph obtained from

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<sup>&</sup>lt;sup>†</sup> Corresponding author: mameij@math.sdu.edu.cn

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the hypercube  $Q_n$  by adding an edge between any two complementary vertices  $x = (x_1 x_2 \cdots x_n)$  and  $\bar{x} = (\bar{x}_1, \bar{x}_2, \cdots, \bar{x}_n)$ , where  $\bar{x}_i = 1 - x_i$ . We call these added edges complementary edges, to distinguish them from the edges, called regular edges, in  $Q_n$ .

From definitions,  $Q_n$  is a proper spanning subgraph of  $FQ_n$ , and so  $FQ_n$  has  $2^n$  vertices. It has been shown that  $FQ_n$  is (n + 1)-regular (n + 1)-connected, and has diameter  $\lceil \frac{n}{2} \rceil$ , about half the diameter of  $Q_n$  [3]. Thus, the folded hypercube  $FQ_n$  is an enhancement on the hypercube  $Q_n$  and has recently attracted many researchers' attention [2, 4, 5, 7, 10, 12]. In this paper, we further investigate other topological properties of  $FQ_n$ , transitivity and panconnectivity.

A graph G is called to be vertex-transitive if for any  $x, y \in V(G)$  there is some  $\sigma \in \text{Aut}(G)$ , the automorphism group of G, such that  $\sigma(x) = y$ ; and edge-transitive if for any  $xy, uv \in E(G)$  there is some  $\phi \in \text{Aut}(G)$ such that  $\{\phi(x), \phi(y)\} = \{u, v\}$ . It has been known that  $Q_n$  is both vertextransitive and edge-transitive [1]. However, the transitivity of  $FQ_n$  has not been proved straightforwardly in the literature. In this paper, we will study some algebraic properties of  $FQ_n$ . Using these properties, we give another proof of a known result that  $FQ_n$  is vertex and edge-transitive.

A graph G is panconnected if for any two different vertices x and y in G and any integer l with  $d_G(x, y) \leq l \leq |V(G)| - 1$  there exists an xy-path of length l, where  $d_G(x, y)$  is the distance between x and y in G [11]. It is easy to see that any bipartite graph with at least three vertices is not panconnected. For this reason, Li *et al* [6] suggested the concept of bipanconnected bipartite graphs. A bipartite graph G is called to be bipanconnected if for any two different vertices x and y in G and any integer l with  $d_G(x, y) \leq l \leq |V(G)| - 1$  such that l and  $d_G(x, y)$  have the same parity there exists an xy-path of length l. Li *et al* [6] have shown that  $Q_n$  is bipanconnected. In this paper, we show that for any two vertices x and y in  $FQ_n$  with distance d,  $FQ_n$  contains an xy-path of length l with  $h \leq l \leq 2^n - 1$  such that l and h have the same parity, where  $h \in$  $\{d, n + 1 - d\}$ . Hence,  $FQ_n$  is bipanconnected if n is odd.

The proofs of our results are in Section 2 and Section3, respectively.

### 2 Algebraic Properties

In this section, we study some algebraic properties of  $FQ_n$ , and as applications, show that  $FQ_n$  is vertex and edge-transitive.

The following notations will be used in the proofs of our main results. The symbol H(x, y) denotes the Hamming distance between two vertices x and y in  $Q_n$ , that is, the number of different bits in the corresponding strings of both vertices. Clearly,  $H(x, y) = d_{Q_n}(x, y)$ . It is also clear that

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 $d_{FQ_n}(x,y) = i$  if and only if H(x,y) = i or n+1-i. Let x = 0u and y = 1v be two vertices in  $FQ_n$ . It is easy to count that

$$H(0u, 1\bar{v}) = H(0u, 1u) + H(1u, 1\bar{v})$$
  
= 1 + [(n - 1) - H(1u, 1v)]  
= n + 1 - H(0u, 1v). (1)

Let  $\Gamma$  be a non-trivial finite group, S be a non-empty subset of  $\Gamma$  without the identity of  $\Gamma$  and with  $S^{-1} = S$ . The Cayley graph  $C_{\Gamma}(S)$  of  $\Gamma$  with respect to S is defined as follows.

$$V = \Gamma;$$
  $(x, y) \in E \Leftrightarrow x^{-1}y \in S, \text{ for any } x, y \in \Gamma.$ 

It has been proved that any Cayley graph is vertex-transitive (see, for example, Theorem 2.2.15 in [14]).

As we have known that the hypercube  $Q_n$  is the Cayley graph  $C_{Z_2^n}(S)$ , where  $Z_2$  denotes the additive group of residue classes modulo 2 on the set  $\{0,1\}, Z_2^n = Z_2 \times Z_2 \times \cdots \times Z_2$ , and  $S = \{(10\cdots 0), (010\cdots 0), \cdots, (0\cdots 010\cdots 0), \cdots, (0\cdots 01)\}$  (see, for example, Example 2 in p89 in [14]). The following theorem shows that  $FQ_n$  is also a Cayley graph.

**Theorem 2.1** The folded hypercube  $FQ_n \cong C_{\mathbb{Z}_2^n}(S \cup \{(11 \cdots 1)\}).$ 

**Proof** Clearly,  $V(FQ_n) = Z_2^n$ . Define a natural mapping

$$\varphi: \quad V(FQ_n) \quad \to \quad Z_2^n \\ x \quad \mapsto \varphi(x) = x.$$

Let x and y be any two vertices in  $FQ_n$ . Since  $(x, y) \in E(FQ_n)$  if and only if H(x, y) = 1 or n. Note that  $x^{-1} = x$  for any  $x \in Z_2^n$ . It follows that H(x, y) = 1 if and only if  $x^{-1}y \in S$ ; and H(x, y) = n if and only if  $x^{-1}y =$  $(11\cdots 1)$ , whereby  $(x, y) \in E(C_{Z_2^n}(S \cup \{(11\cdots 1)\}))$ . Thus,  $\varphi$  preserves the adjacency of vertices, which implies that  $\varphi$  is an isomorphism between  $FQ_n$  and  $C_{Z_2^n}(S \cup \{(11\cdots 1)\})$ , and so  $FQ_n \cong C_{Z_2^n}(S \cup \{(11\cdots 1)\})$ .

**Corollary 2.2** The folded hypercube  $FQ_n$  is vertex-transitive.

For convenience, we express  $FQ_n$  as  $FQ_n = L \otimes R$ , where L and R are the two (n-1)-dimensional subcubes of  $Q_n$  induced by the vertices with the leftmost bit is 0 and 1, respectively. A vertex in L will be denoted by 0u and a vertex in R denoted by 1v, where u and v are any two vertices in  $Q_{n-1}$ . Between L and R, apart from the regular edges, there exists a complementary edge joining 0u and  $1\bar{u} \in R$  for any  $0u \in L$ .

**Theorem 2.3** Let  $\sigma$  be a mapping from  $V(FQ_n)$  to itself defined by

$$\begin{cases} \sigma(0u) = 0u \\ \sigma(1u) = 1\bar{u} \end{cases} \quad \text{for any } u \in V(Q_{n-1}).$$

$$\tag{2}$$

Then  $\sigma \in \text{Aut}(FQ_n)$ . Moreover, for an edge (x, y) between L and R in  $FQ_n$ ,  $(\sigma(x), \sigma(y))$  is complementary if and only if (x, y) is regular.

**Proof** Clearly,  $\sigma$  is a permutation on  $V(FQ_n)$ . To show  $\sigma \in \text{Aut}(FQ_n)$ , it is sufficient to show that  $\sigma$  preserves adjacency of vertices in  $FQ_n$ , that is, to show that any pair of vertices x and y in  $FQ_n$  satisfies the following condition.

$$(x,y) \in E(FQ_n) \Leftrightarrow (\sigma(x), \sigma(y)) \in E(FQ_n).$$
(3)

Let  $FQ_n = L \otimes R$ , u and v be any two distinct vertices in  $Q_{n-1}$ . Because of vertex-transitivity of  $FQ_n$  by Theorem 2.1, without loss of generality, suppose  $x = 0u \in L$ . We consider two cases according to the location of y.

Case 1  $y \in L$ . In this case, let y = 0v. Since  $\sigma$  is the identical permutation on  $L \cong Q_{n-1}$ , it is clear that

$$(0u, 0v) \in E(FQ_n) \Leftrightarrow (\sigma(0u), \sigma(0v)) = (0u, 0v) \in E(FQ_n).$$

Case 2  $y \in R$ . In this case, let y = 1v. By the definition of  $FQ_n$ ,  $(0u, 1v) \in E(FQ_n) \Leftrightarrow v = u$  or  $\bar{u}$ . Since  $(0u, 1u), (0u, 1\bar{u}) \in E(FQ_n)$  by the definition of  $FQ_n$ , it follows that

$$(0u, 1u) \in E(FQ_n) \Leftrightarrow (\sigma(0u), \sigma(1u)) = (0u, 1\bar{u}) \in E(FQ_n) \quad \text{if } v = u \\ (0u, 1\bar{u}) \in E(FQ_n) \Leftrightarrow (\sigma(0u), \sigma(1\bar{u})) = (0u, 1u) \in E(FQ_n) \quad \text{if } v = \bar{u}.$$

From the above arguments, we have shown  $\sigma \in \operatorname{Aut}(FQ_n)$ .

We now show the remaining part of the theorem. Without loss of generality, we may suppose x = 0u since  $FQ_n$  is vertex-transitive. By (2), we have  $\sigma(x) = \sigma(0u) = 0u$ .

Suppose that (x, y) is a regular edge between L and R in  $FQ_n$ . Then y = 1u and  $\sigma(y) = \sigma(1u) = 1\overline{u}$  by (2). By (3)  $(0u, 1\overline{u}) \in E(FQ_n)$ , which is a complementary edge.

Conversely, suppose that (x, y) is a complementary edge in  $FQ_n$ . Then  $y = 1\overline{u}$  and  $\sigma(1\overline{u}) = 1u$  by (2), and  $(0u, 1u) \in E(FQ_n)$  by (3), which is a regular edge.

The lemma follows.

**Theorem 2.4** Aut  $(Q_n)$  is a proper subgroup of Aut  $(FQ_n)$ . Moreover, for any  $\sigma \in \text{Aut}(Q_n)$ , (x, y) is a complementary edge if and only if  $(\sigma(x), \sigma(y))$  is also a complementary edge in  $FQ_n$ .

**Proof** For any element  $\sigma \in \operatorname{Aut}(Q_n)$ , we will prove  $\sigma \in \operatorname{Aut}(FQ_n)$ .

It is clear that  $\sigma$  is a permutation on  $V(FQ_n)$  since  $Q_n$  is a spanning subgraph of  $FQ_n$ . We only need to show that  $\sigma$  preserves adjacency of vertices in  $FQ_n$ , that is, to check that (3) holds for any pair of vertices xand y in  $FQ_n$ . In fact, since

$$H(x,y) = d_{Q_n}(x,y) = d_{Q_n}(\sigma(x),\sigma(y)) = H(\sigma(x),\sigma(y))$$

$$(x, y) \in E(FQ_n) \Leftrightarrow H(x, y) = 1 \text{ or } n,$$

we have

$$\begin{array}{ll} (x,y) \in E(FQ_n) & \Leftrightarrow & H(x,y) = 1 \text{ or } n \\ & \Leftrightarrow & H(\sigma(x),\sigma(y)) = 1 \text{ or } n \\ & \Leftrightarrow & (\sigma(x),\sigma(y)) \in E(FQ_n). \end{array}$$

Thus, Aut  $(Q_n) \subseteq$  Aut  $(FQ_n)$ . It is clear that the automorphism  $\sigma$  defined by (2) is not in Aut  $(Q_n)$  by Theorem 2.3. Therefore, Aut  $(Q_n)$  is a proper subgraph of Aut  $(FQ_n)$ .

By the definition of  $FQ_n$ , for any  $\sigma \in \operatorname{Aut}(Q_n)$ , it is clear that (x, y) is a complementary edge in  $FQ_n$  if and only if  $n = H(x, y) = H(\sigma(x), \sigma(y))$ , if and only if  $\sigma(x, y) = (\sigma(x), \sigma(y))$  is a complementary edge in  $FQ_n$ .

The theorem follows.

**Corollary 2.5** The folded hypercube  $FQ_n$  is edge-transitive.

**Proof** For any two edges (x, y) and (x', y') in  $FQ_n$ , we will show there is an element  $\sigma \in \operatorname{Aut}(FQ_n)$  such that  $\{\sigma(x), \sigma(y)\} = \{x', y'\}$ . Since  $FQ_n$  is vertex-transitive, we may assume x = x'. We only need to find  $\sigma \in \operatorname{Aut}(FQ_n)$  that takes y to y' and fixes x. Since for any two vertices zand t in  $FQ_n$ ,  $(z, t) \in E(FQ_n)$  if and only if H(z, t) = 1 or n. Without loss of generality, we may suppose that H(x, y) = 1, that is, (x, y) is a regular edge in  $FQ_n$ .

If H(x, y') = 1, then (x, y') is a regular edge. Since  $Q_n$  is edge-transitive, there is an element  $\sigma \in \text{Aut}(Q_n)$  such that  $\{\sigma(x), \sigma(y)\} = \{x, y'\}$ . By Theorem 2.4,  $\sigma \in \text{Aut}(FQ_n)$ , which satisfies our requirement.

If H(x, y') = n, then  $y' = \bar{x}$  and (x, y') is a complementary edge in  $FQ_n$ . Without loss of generality, we may suppose that x = 0u. Then  $y' = 1\bar{u}$ . Let z = 1u. Then the automorphism  $\sigma$  defined in (2) can take z to y' and fixes x. If y = z, then the  $\sigma$  satisfies our requirement. If  $y \neq z$ , then there is  $\phi \in \operatorname{Aut}(Q_n) \subset \operatorname{Aut}(FQ_n)$  such that  $\phi$  takes y to z and fixes x. Thus,  $\sigma\phi(y) = \sigma(\phi(y)) = \sigma(z) = y'$  and  $\sigma\phi(x) = \sigma(\phi(x)) = \sigma(x) = x$ , and so  $\sigma\phi$  satisfies our requirement.

The corollary follows.

## **3** Panconnectivity

In this section, we investigate the panconnectivity of  $FQ_n$ . The proof of the main theorem in this section is strongly dependent on the following lemmas.

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and

**Lemma 3.1** [6] If  $n \ge 2$ , then  $Q_n$  is bipanconnected, that is, for any two vertices x and y in  $Q_n$  there exists an xy-path of length l with  $H(x, y) \le l \le 2^n - 1$  such that l and H(x, y) have the same parity.

**Lemma 3.2** [13]  $FQ_n$  is a bipartite graph if and only if n is odd. Moreover, if n is even, then the length of the shortest odd cycle in  $FQ_n$  is n+1.

**Theorem 3.3** For any two distinct vertices x and y in  $FQ_n$  with distance d,  $FQ_n$  contains an xy-path of length l with  $h \le l \le 2^n - 1$  such that l and h have the same parity, where  $h \in \{d, n + 1 - d\}$ .

**Proof** If n = 1, the theorem is true clearly since  $FQ_1 = K_2$ . Assume  $n \ge 2$  below. Without loss of generality, we may assume x = 0u, y = 1v since  $d \ge 1$  and  $FQ_n$  is vertex-transitive by Corollary 2.2. We first deduce two conclusions from Lemma 3.2 and Theorem 2.3.

(a) By Lemma 3.1,  $Q_n$  contains an xy-path P of length l with  $H(x, y) \leq l \leq 2^n - 1$  such that l and H(x, y) have the same parity. Since  $Q_n$  is a spanning subgraph of  $FQ_n$ , P is an xy-path of length l in  $FQ_n$ .

(b) Consider the vertex  $z = 1\overline{v}$ . By Lemma 3.1,  $Q_n$  contains an xz-path R of length l' with  $H(x, z) \leq l' \leq 2^n - 1$  such that l' and H(x, z) have the same parity. Since H(x, z) = n + 1 - H(x, y) by (1), l' and n + 1 - H(x, y) have the same parity. Let  $\sigma \in \text{Aut}(FQ_n)$  defined in (2). Then  $P' = \sigma(R)$  is an xy-path of length l' with  $n + 1 - H(x, y) \leq l' \leq 2^n - 1$  such that l' and n + 1 - H(x, y) have the same parity.

To prove the theorem, it is sufficient to check that H(x,y) = d or n+1-d. In fact, it is clear that if  $H(x,y) \leq \lfloor \frac{n}{2} \rfloor$  then d = H(x,y); if  $H(x,y) > \lfloor \frac{n}{2} \rfloor$  then H(x,y) = n-d+1. The theorem is proved.

**Corollary 3.4** If n is odd, then  $FQ_n$  is bipanconnected.

**Proof** If *n* is odd, then  $FQ_n$  is a bipartite graph by Lemma 3.2. Let x and y be any two vertices in  $FQ_n$  with distance d. Since *n* is odd, the condition that l and n + 1 - d have the same parity implies that l and d have the same parity. Note that  $d \le n + 1 - d$  since  $d \le \lfloor \frac{n}{2} \rfloor$ . By Theorem 3.3,  $FQ_n$  contains an xy-path of length l with  $d \le l \le 2^n - 1$  such that l and d have the same parity, and so  $FQ_n$  is bipanconnected.

**Corollary 3.5** If *n* is even then for any two different vertices *x* and *y* with  $d_{FQ_n}(x, y) = d$  in  $FQ_n$ , there is an *xy*-path of length *l* for each *l* satisfying  $n - d + 1 \le l \le 2^n - 1$  and there is also an *xy*-path of length *l'* for each *l'* satisfying  $d \le l' \le n - d$  such that *l'* and *d* have the same parity; there is no *xy*-path of other length.

**Proof** If n is even, then d and n - d + 1 have different parity. Thus, for any integer l, either l and d have the same parity, or l and n - d + 1

have the same parity. Since  $d \leq \frac{n}{2}$ , d < n - d + 1. By Theorem 3.3, there is an *xy*-path of length l with  $n - d + 1 \leq l \leq 2^n - 1$  in  $FQ_n$ .

Since the length of the shortest odd cycle in  $FQ_n$  is n + 1 by Lemma 3.2,  $FQ_n$  contains no xy-path of length l with  $d < l \leq n - d$  if l and d have different parity. In other words, the length l of the second shortest path between x and y with distance d is certainly n - d + 1 if l and d have different parity. It follows from Theorem 3.3 that there is an xy-path of length l' with  $d \leq l' \leq n - d$  provided l' and d have the same parity.

The corollary is proved.

A graph is called to be hamiltonian connected if there is a hamiltonian path between any two vertices. It is easy to see that any bipartite graph with at least three vertices is not hamiltonian connected. For this reason, Simmons [9] introduces the concept of hamiltonian laceable for hamiltonian bipartite graphs. A hamiltonian bipartite graph is hamiltonian laceable if there is a hamiltonian path between any two vertices in different bipartite sets. It is clear that if a bipartite graph is bipanconnected then it is certainly hamiltonian laceable. It follows from Corollary 3.4 and Corollary 3.5 that the following result is true clearly.

**Corollary 3.6**  $FQ_n$  is hamiltonian laceable if n is odd, and hamiltonian connected if n is even.

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